Dense Linear Algebra Solvers for Multicore with GPU Accelerators

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IEEE IPDPS 2010
High-level Parallel Programming Models and Supportive Environments (HIPS)
April 19-23, 2010, Atlanta, GA



Outline

Introduction

Hardware to Software Trends

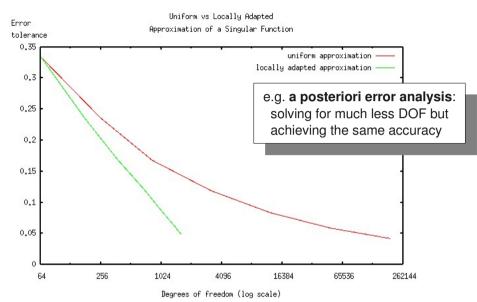
The MAGMA library

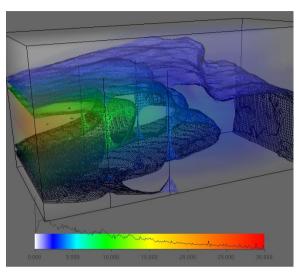
- Challenges and approach
- One-sided factorizations and solvers
- Two-sided factorizations

Conclusions

Speeding up Computer Simulations

Better numerical methods

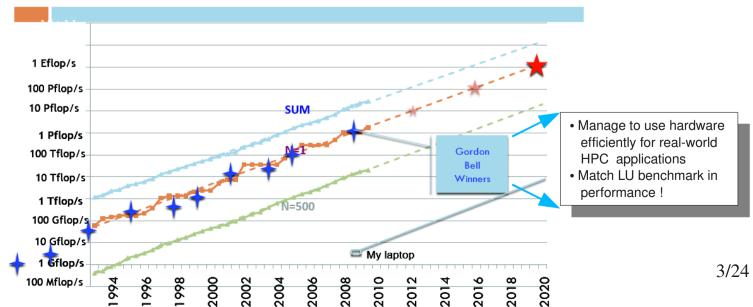




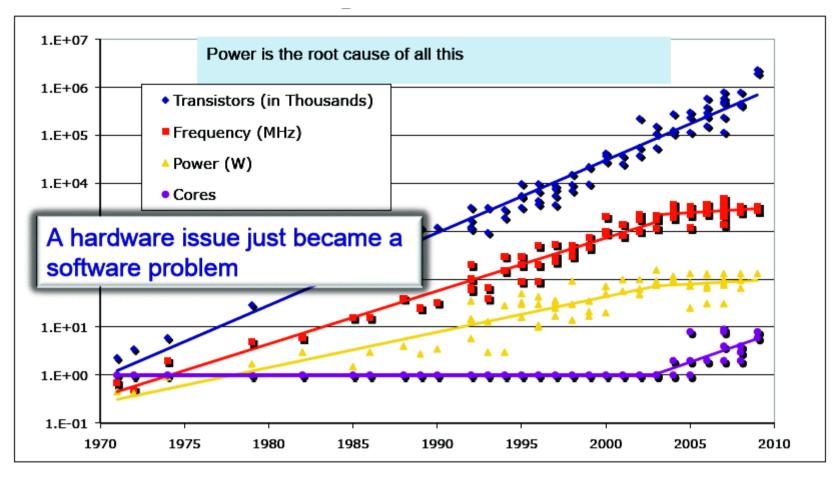
http://www.cs.utk.edu/~tomov/cflow/

Performance Development in Top500

Exploit advances in hardware



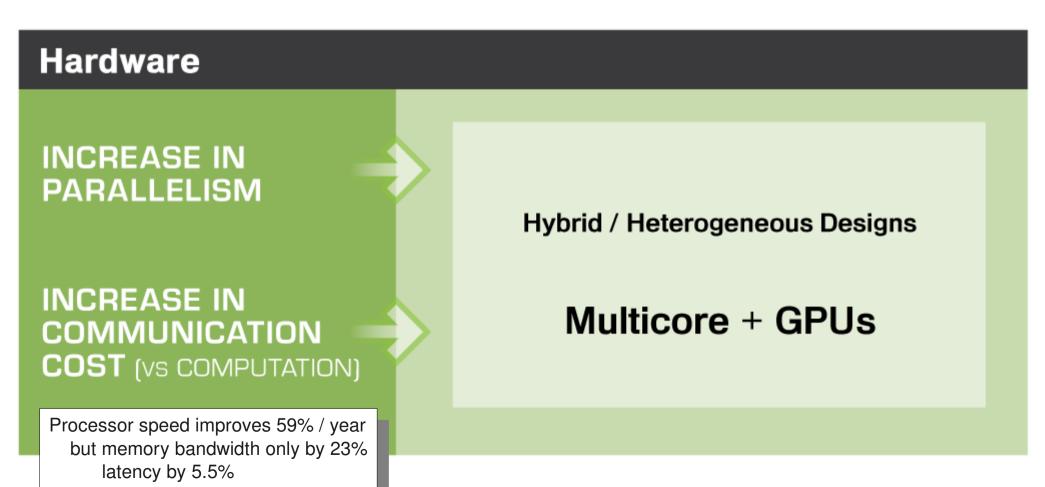
Clock Frequency Scaling Replaced by Scaling Cores/Chip



Data from Kunle Olukotun, Lance Hammond, Herb Sutter, Burton Smith, Chris Batten, and Krste Asanoviç Slide from Kathy Yelick

Why GPU-based Computing?

Hardware Trends



Matrix Algebra on GPU and Multicore Architectures (MAGMA)

MAGMA: a new generation linear algebra (LA) libraries to achieve the fastest possible time to an accurate solution on hybrid/heterogeneous architectures, starting with current multicore+MultiGPU systems Homepage: http://icl.cs.utk.edu/magma/

MAGMA & LAPACK

- MAGMA based on LAPACK and extended for hybrid systems (multi-GPUs + multicore systems);
- MAGMA designed to be similar to LAPACK in functionality, data storage and interface, in order to allow scientists to effortlessly port any of their LAPACK-relying software components to take advantage of the new architectures
- MAGMA to leverage years of experience in developing open source LA software packages and systems like LAPACK, ScaLAPACK, BLAS, ATLAS as well as the newest LA developments (e.g. communication avoiding algorithms) and experiences on homogeneous multicores (e.g. PLASMA)

Support

 NSF, Microsoft, NVIDIA [now CUDA Center of Excellence at UTK on the development of Linear Algebra Libraries for CUDA-based Hybrid Architectures]

MAGMA developers

University of Tennessee, Knoxville; University of California, Berkeley; University of Colorado, Denver

MAGMA 0.2

- LU, QR, Cholesky (S, C, D, Z)
- Linear solvers
 - In working precision, based on LU, QR, and Cholesky
 - Mixed-precision iterative refinement
- CPU and GPU interfaces
- Two-sided factorizations
 - Reduction to upper Hessenberg form (bi/tri-diagonalization developed)
- MAGMA BLAS
 - Routines critical for MAGMA (GEMM, SYRK, TRSM, GEMV, SYMV, etc.)

Challenges

Massive parallelism

Many GPU cores, serial kernel execution [e.g. 240 in the GTX280; up to 512 in *Fermi* – to have concurrent kernel execution]

Hybrid/heterogeneous architectures

Match algorithmic requirements to architectural strengths [e.g. small, non-parallelizable tasks to run on CPU, large and parallelizable on GPU]

Compute vs communication gap

Exponentially growing gap; persistent challenge

[on all levels, e.g. a GPU Tesla C1070 (4 x C1060) has compute power of O(1,000) Gflop/s but GPUs communicate through the CPU using O(1) GB/s connection]

How to Code for GPUs?

Complex question

- Language, programming model, user productivity, etc
- Recommendations
 - Use CUDA / OpenCL

[already demonstrated benefits in many areas; data-based parallelism; move to support task-based]

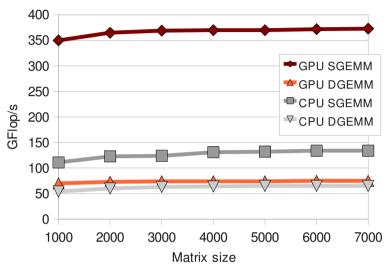
Use GPU BLAS

[high level; available after introduction of **shared memory** – can do data reuse; leverage existing developments]

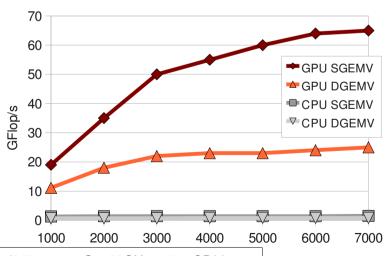
Use Hybrid Algorithms

[currently GPUs – massive parallelism but serial kernel execution; hybrid approach – small non-parallelizable tasks on the CPU, large parallelizable tasks on the GPU]





GPU vs CPU GEMV



GPU: GTX280 (240 cores @ 1.30GHz, 141 GB/s) **CPU**: 2 x 4 cores Intel Xeon @ 2.33GHz, 10.4 GB/s)

9/24

LAPACK to Multicore

A New Generation of Software:

Parallel Linear Algebra Software for Multicore Architectures (PLASMA)

Software/Algorithms follow hardware evolution in time		
LINPACK (70's) (Vector operations)		Rely on - Level-1 BLAS operations
LAPACK (80's) (Blocking, cache friendly)		Rely on - Level-3 BLAS operations
ScaLAPACK (90's) (Distributed Memory)		Rely on - PBLAS Mess Passing
PLASMA (00's) New Algorithms (many-core friendly) Those new algorithms		Rely on - a DAG/scheduler - block data layout - some extra kernels

"delayed update" to organize successive Level 2 BLAS as a single Level 3 BLAS

Split BLAS into tasks and represent algorithms as DAGs; new algorithms where panel factorizations use localized (over tiles) elementary transformations

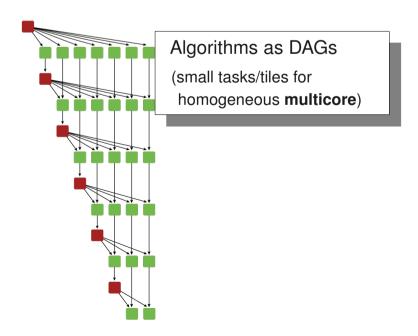
- have a very low granularity, they scale very well (multicore, petascale computing, ...)
- removes a lots of dependencies among the tasks, (multicore, distributed computing)
- avoid latency (distributed computing, out-of-core)
- rely on fast kernels

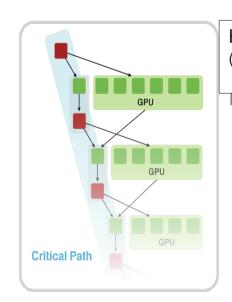
Those new algorithms need new kernels and rely on efficient scheduling algorithms.

LAPACK to MAGMA

(multicore with GPU accelerators)

- 1) Development of NEW LGORITHMS (parallelism, hybrid, optimized communication)
- 2) HYBRIDIZATION of linear algebra algorithms
 - Represent the algorithms as a collection of TASKS and DEPENDANCIES among them
 - Properly SCHEDULE the tasks' execution over the multicore and the GPU
- 3) Development of GPU BLAS **KERNELS**
- 4) **AUTO-TUNED** implementations

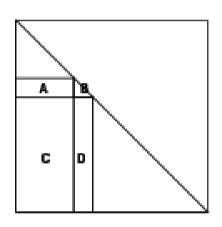




Hybrid CPU+GPU algorithms (small tasks for multicores and large tasks for GPUs)

One-Sided Dense Matrix Factorizations (LU, QR, and Cholesky)

- Panels (Level 2 BLAS) are factored on CPU using LAPACK
- Trailing matrix updates (Level 3 BLAS) are done on the GPU using "look-ahead" (to overlap CPUs work on the critical path with the GPUs large updates)



Example: Left-Looking Hybrid Cholesky factorization

MATLAB code

(1) B = B - A*A

$$(4) D = D \setminus B$$

LAPACK code

```
ssyrk_("L", "N", 8mb, 8j, 8mone, hA(j,0), ... ]
spotrf_("L", 8mb, hA(j,j), Ide, info)
sgomm_("N", "T", 8j,...)
```

strsm ("R", "L", "T", "N", Bi, ...]

Hybrid code

```
cublasSsyrk['L', 'N', nb, j. mone, dA[j,0], ...]

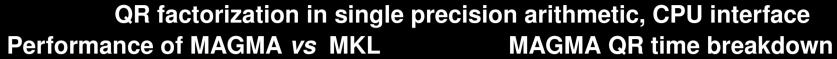
cublasGetMatrix[nb, nb, 4, dA[j, j], *lda, hwork, nb]

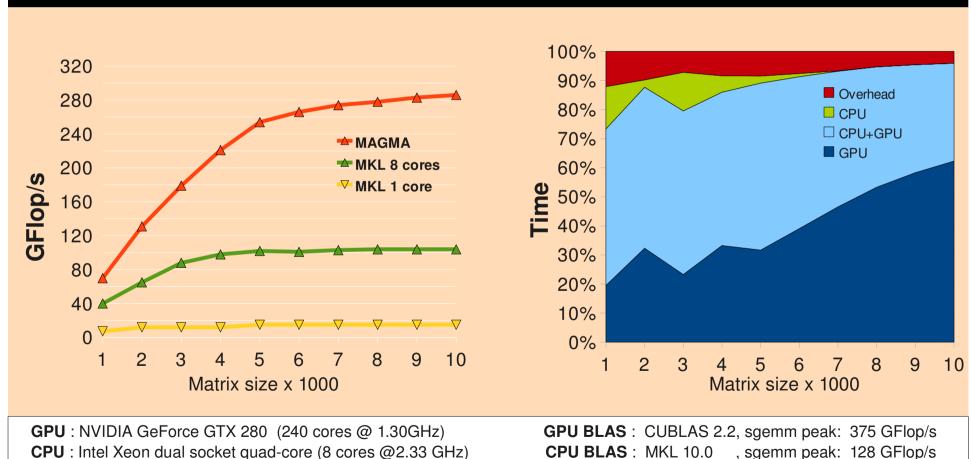
cublasSgemm['N', 'T', j, ...]

spotrf_['L', 8mb, hwork, 8mb, info]

cublasSetMatrix[nb, nb, 4, hwork, nb, dA[j, j], *lda]
```

One-sided hybrid factorizations

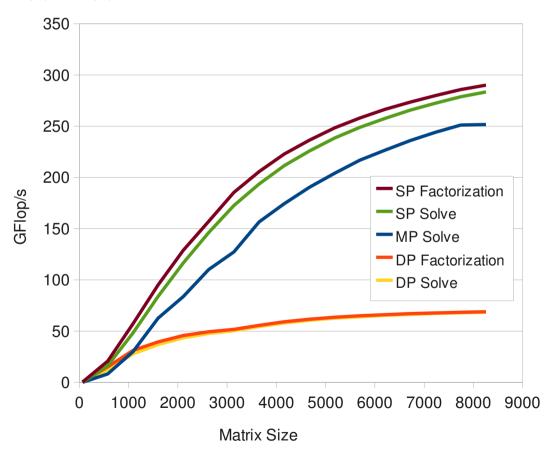




Linear Solvers

Solving Ax = b using LU factorization

Intel(R) Xeon(R)E541@2.34GHz / 8 Cores + GTX 280 @1.30GHz / 240 Cores

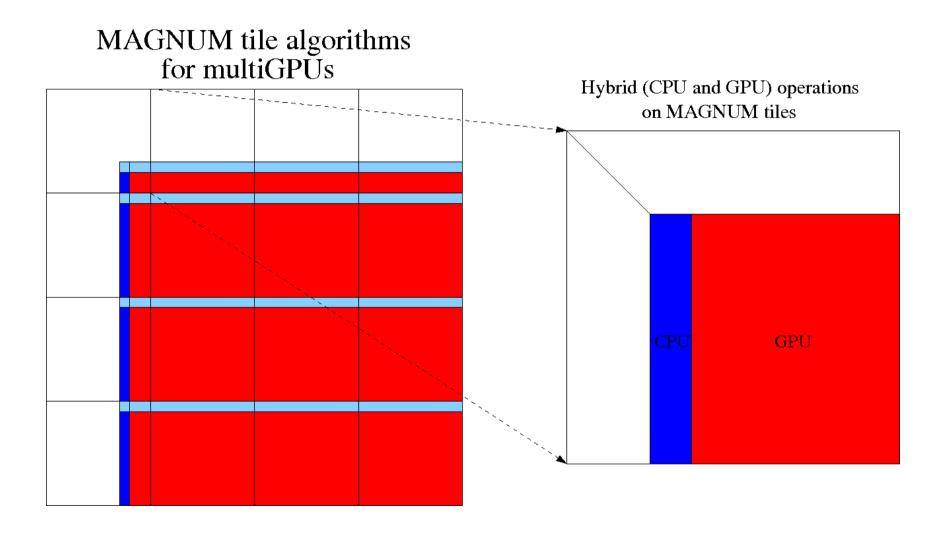


Direct solvers

- Factor and do triangular solves in the same, working precision
- Mixed Precision Iterative Refinement
 - Factor in single (i.e. the bulk of the computation in fast arithmetic) and use it as preconditioner in simple double precision iteration, e.g.

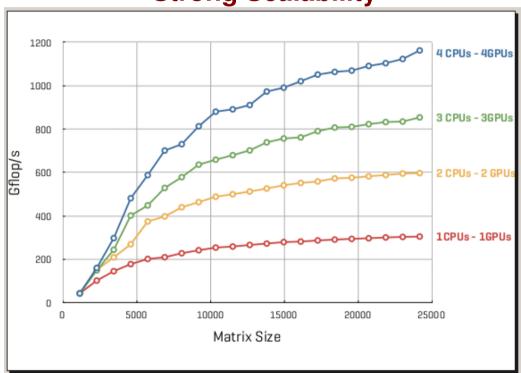
$$X_{i+1} = X_i + (LU_{SP})^{-1} P (b - A X_i)$$

Extension to Multicore and Multi GPUs



Performance using MultiGPUs

Cholesky factorization in SP Strong Scalability



HOST: 4x AMD Opteron core @1.8GHz

GPUs: 4x C1060 (240 cores each @1.44GHz)

2 level nested parallelism

coarse: PLASMA tiled algorithm and

static scheduling

fine : tasks/tiles are redefined for

hybrid 1 core+GPU computing

- Defining a "Magnum tiles approach"

Two-sided matrix factorizations

Two-sided factorizations

$$QAQ'=H$$

H – upper Hessenberg / bidiagonal / tridiagonal,

Q – orthogonal similarity transformation

Importance

One-sided factorizations

- bases for linear solvers

Two-sided factorizations

- bases for eigen-solvers

Block algorithm

Q – a product of n-1 elementary reflectors

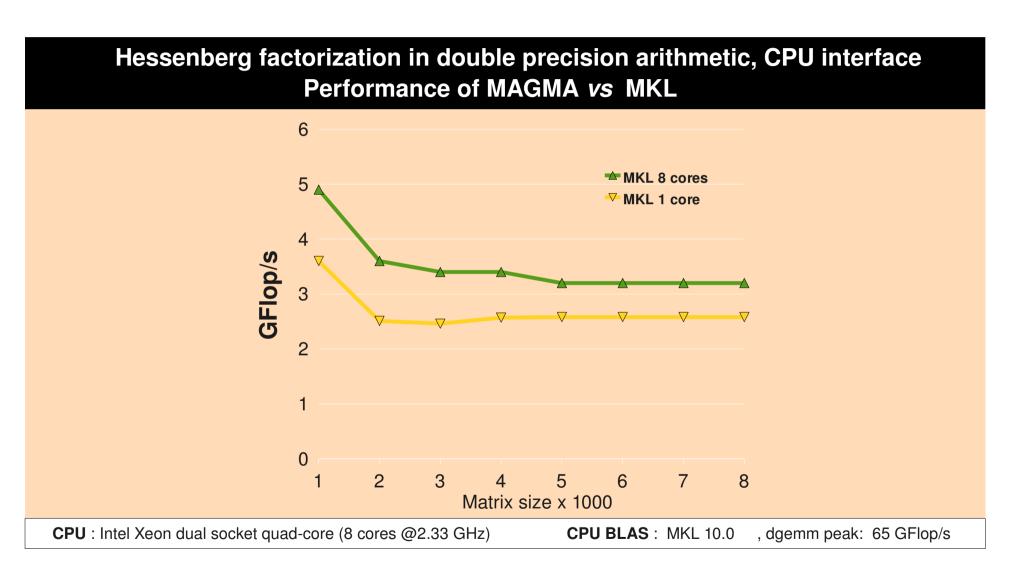
$$Q = H_1 H_2 ... H_{n-1}, H_i = I - \tau_i v_i v_i'$$

 $H_1 \dots H_{nb} = I - V T V'$ (WY transform; the bases for delayed update or block algorithm)

Can we accelerate it ?

[similarly to the one-sided using hybrid GPU-based computing]
[to see **much higher acceleration** due to a removed bottleneck] 17/24

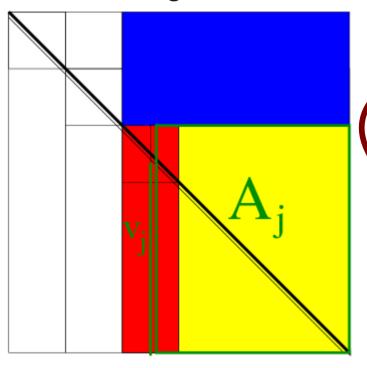
Homogeneous multicore acceleration?



There have been difficulties in accelerating it on homogeneous multicores

The Bottleneck

Hessenberg factorization



■ Level 3 BLAS update

[80% flops; ~30% of the run time]

■ Level 2 BLAS update [20% flops; ~70% of the run time]

 $y_{j} = A_{j} v_{j}$

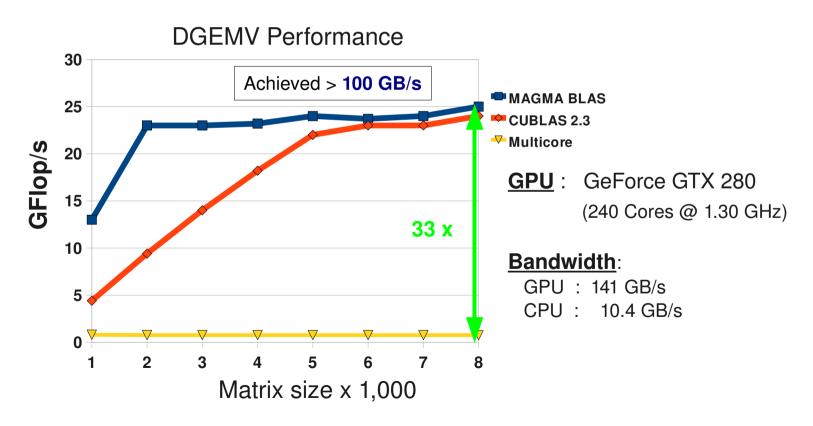
bidiagonalization & tridiagonalization have even more Level 2 BLAS (50%)

Reduction times in seconds for N = 4,000

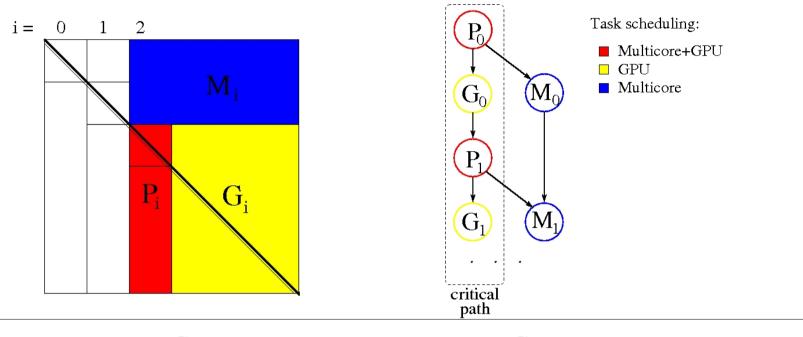


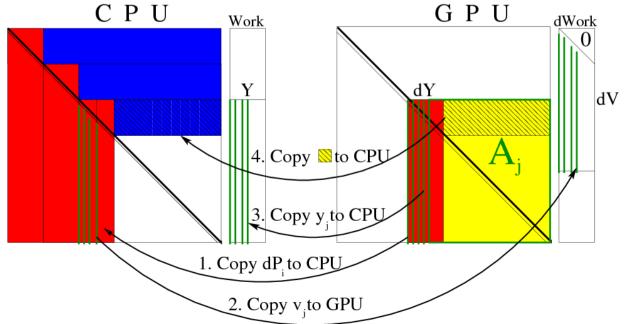
Hybrid computing acceleration?

- Intuitively, yes, as matrix-vector product is fast on GPUs (e.g., sgemv up to 66 Gflop/s, ssymv up to 102 GFlop/s)
- How to organize a hybrid computation?



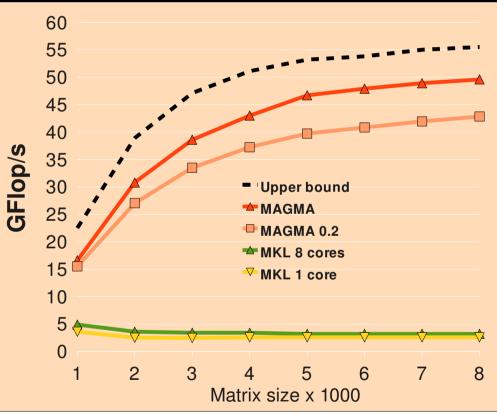
Task Splitting & Task Scheduling





Performance

Hessenberg factorization in double precision arithmetic, CPU interface Performance of MAGMA vs MKL



GPU: NVIDIA GeForce GTX 280 (240 cores @ 1.30GHz)

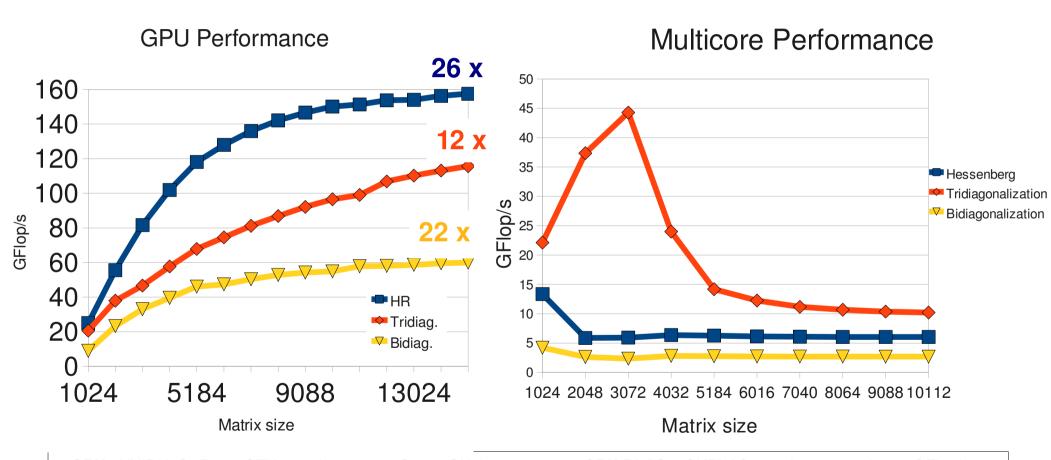
CPU: Intel Xeon dual socket quad-core (8 cores @2.33 GHz)

GPU BLAS: CUBLAS 2.3, dgemm peak: 75 GFlop/s

CPU BLAS: MKL 10.0 , dgemm peak: 65 GFlop/s

Two-sided factorizations

(performance in single precision arithmetic)



GPU: NVIDIA GeForce GTX 280 (240 cores @ 1.30GHz) **CPU**: Intel Xeon dual socket quad-core (8 cores @2.33 GHz)

GPU BLAS: CUBLAS 2.3, dgemm peak: 75 GFlop/s **CPU BLAS**: MKL 10.0 , dgemm peak: 65 GFlop/s

Conclusions

- Linear algebra can be significantly accelerated using GPUs
- Described a hybridization methodology to achieve this acceleration
 - high level model
 - Leverage prior developments
- Hybridization can be used for a wide set of fundamental linear algebra algorithms
 - Linear and eigen/singular-value solvers
 - Incorporated in the MAGMA library http://icl.cs.utk.edu/magma/