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KEENELAND: BRINGING HETEROGENEOUS GPU COMPUTING TO THE COMPUTATIONAL SCIENCE COMMUNITY

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The Keeneland project's goal is to develop and deploy an innovative, GPU-based high-performance computing system for the NSF computational science community.

he Keeneland project—named for a historic thoroughbred horse racing track in Lexington, Kentucky—is a five-year Track 2D grant awarded by the US National Science Foundation (NSF) in August 2009 for the development and deployment of an innovative high-performance computing system. The Keeneland project is led by the Georgia Institute of Technology (Georgia Tech) in collaboration with the University of Tennessee at Knoxville and Oak Ridge National Laboratory. The initial delivery system, an HP Linux cluster with 360 Nvidia Fermi graphics processors, is now operational and is being used to prepare software tools and applications for the full-scale system, which is due to be deployed in 2012.

Keeneland is organized into two primary deployment phases. The first phase provides a moderately sized, initial delivery system (KID) to develop software tools for GPU computing and to prepare applications to exploit GPUs effectively. During 2012, Keeneland's second phase will provide a full-scale system for production use by computational scientists. The Keeneland Full Scale (KFS) system will be similar to the KID system in terms of hardware and software—that is, it will be a Linux cluster based on

commodity next-generation CPUs, next-generation GPUs, HPC interconnect, programming environment, and tools. The KFS system is expected to have at least two times KID's total number of GPUs, with the expected performance improvements of the next-generation CPUs and GPUs. KFS will be an Extreme Digital Science and Engineering Discovery Environment (XSEDE) resource available to a broad set of users. Although there's currently a community movement in HPC toward this type of architecture, a critical component of the Keeneland project is to develop software that will let users take advantage of its unique capabilities. We also aim to reach out to teams developing applications that might map well to this innovative architecture.

Motivation

Heterogeneous architectures have recently emerged in response to the limits on improving single-node performance in traditional architectures; the primary factor influencing these limits is energy efficiency. In the case of GPUs,^{1,2} these heterogeneous architectures were initially designed to support a fixed pipeline of graphics operations (such as rasterization), and consequently, architects

could specialize them to be exceedingly efficient at those operations. However, it was only after a small set of early adopters began using GPUs for general-purpose computation more than a decade ago^{3,4} that hardware features, algorithmic techniques, and programming systems—such as Cg,⁵ CUDA,⁶ and OpenCL^{7,8}—started to emerge to make GPUs available to a wider audience.

Most recently, GPUs such as Nvidia's Fermi¹ have added critical features, including much-improved performance on IEEE doubleprecision arithmetic and memory error detection and correction, that make these architectures even more relevant to large-scale computational science. Compared to the alternatives, these new features, when combined with the original capabilities of GPUs, provide a competitive platform for numerous types of computing, such as media processing, gaming, and scientific computing, in terms of raw performance (665 Gflops/s per Fermi) and energy efficiency.

Not surprisingly, these trends have garnered the attention of researchers, vendors, and HPC customers. Beyond the Keeneland project, a substantial number of very large GPU-based systems have already been deployed.

Examples include China's Tianhe-1A, Nebulae at the National Supercomputing Centre (NSCS) in Shenzhen, Tokyo Tech's Tsubame 2, and several other systems in the US including Dirac at Lawrence Berkeley National Laboratory, Lincoln at the National Center for Supercomputing Applications (NCSA), and Edge at the Lawrence Livermore National Laboratory. Notably, the Chinese Tianhe-1A system at the NSCS in Tianjin achieves a performance of 2.57 petaflops/s on the TOP500 Linpack benchmark (www.top500.org), which was #1 in the world in November 2010.

All of these examples are scalable heterogeneous architectures that leverage commodity components: multinode computing systems with a highperformance interconnection network, where each node contains more than one type of processing device, such as a traditional CPU and a graphics processor. Most experts expect this trend to continue into the foreseeable future, given the requirements and constraints of HPC. Even with this tremendous progress over the past few years, the adoption of GPUs by the wider computational science community faces several hurdles. These hurdles include programmability; portability; consistent performance; limitations of architectural interfaces; and the fact that some application features, such as performing I/O, simply won't perform well on current GPU architectures.

Keeneland

The Keeneland project's first phase is underway; it includes the acquisition and operation of the initial delivery system, development of software tools, preparation of applications for GPU computing, and an assessment of the fast-changing technologies.

Table 1. The KID Configuration.		
Node architecture	HP ProLiant SL390 G7	
CPU	Intel Xeon X5660 (Westmere)	
CPU frequency	2.80 Ghz	
CPU cores per node	12	
Host memory per node	24 Gbytes	
GPU architecture	Nvidia Tesla M2070 (Fermi)	
GPUs per node	3	
GPU memory per node	18 Gbytes (6 Gbytes per GPU)	
CPU/GPU ratio	2:3	
Interconnect	InfiniBand QDR (single rail)	
Total number of nodes	120	
Total CPU cores	1,440	
Total GPU cores	161,280	

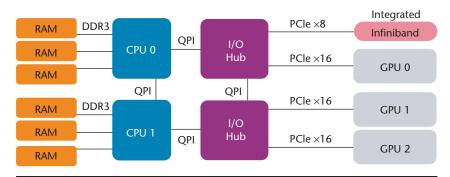


Figure 1. The HP SL390 G7 Node Architecture provides two host CPUs, up to three GPUs, and an integrated Infiniband QDR HCA.

Architecture

KID has been installed and operating since November 2010. As Table 1 shows, KID's configuration is rooted in the scalable node architecture of the HP ProLiant SL390 G7. In particular, as Figure 1 shows, each node has two Intel Westmere CPUs, three Nvidia M2070 Fermi GPUs with 6 Gbytes of memory each, 24 Gbytes of host main memory, and a Mellanox Quad Data Rate (QDR) InfiniBand Host Channel Adapter (HCA). Overall, the system has 120 nodes with 240 CPUs and 360 GPUs; the installed system has a peak performance of 201 Tflops in seven racks (or 90 square feet, including the service area).

More specifically, in the HP SL390, memory is directly attached to the CPU sockets, which are connected to each other and the Tylersburg I/O

hubs via Intel's Quick Path Interconnect (QPI). GPUs are attached to the node's two I/O hubs using Peripheral Component Interconnect Express (PCIe). The theoretical peak for QPI's unidirectional bandwidth is approximately 12.8 Gbytes/s, and for PCIe x16 it is approximately 8.0 Gbytes/s. With these two I/O hubs, each node can simultaneously supply three full x16 PCIe links to the GPUs and an x8 link to the integrated Infiniband QDR HCA.

This design avoids contention and offers advantages in aggregate node bandwidth when the three GPUs and the HCA are used concurrently, as they often are in a scientific system. In contrast, other architectures frequently use a PCIe-switch-based approach, and the switch can quickly become a performance bottleneck.

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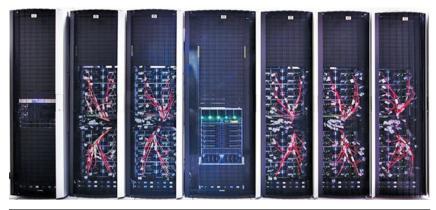


Figure 2. The KID system as installed in Keeneland's data center. The compact 201 teraflops/s system requires only seven racks and 90 square feet of floor space.

Nevertheless, with this PCIe-switchbased approach, vendors are currently offering systems with the highest number of GPUs per node.

The node architecture exemplifies the heterogeneity trends and has one of the highest number of GPU counts per node in the November TOP500 list. The HP SL390 design has significant benefits over the previous generation architecture, but also exhibits multiple levels of non-uniformity.9 In addition to traditional non-uniform memory access (NUMA) effects across the two Westmere's integrated memory controllers (see Figure 1), the dual I/O hub design introduces non-uniform characteristics for data transfers between host memory and GPU memory. These transfers will perform better if the data traverses only one QPI link (such as a transfer between data in the memory attached to CPU socket 0 and GPU 0) than if it traverses two QPI links (such as a transfer between data in the memory attached to CPU socket 0 and GPU 1 or GPU 2).

In addition, KID's GPUs include other features that can greatly affect performance and contribute to non-uniformity. For instance, each GPU contains error-correcting code (ECC) memory. ECC memory is desirable in a system designed for scalable scientific computing. Enabling ECC gives some assurance against these transient errors, but results in a

performance penalty and adds yet another complexity to the GPU memory hierarchy.

Software Tools

As mentioned earlier, one of the risks in using these new heterogeneous platforms involves decreased programmer productivity. On the Keeneland project, we've started collaborating with many vendors and initiated research activities to address this challenge. In particular, software tools currently developed under the Keeneland project include scientific libraries, performance and correctness tools, and virtualization system software.

For scientific libraries, the University of Tennessee at Knoxville's Matrix Algebra on GPU and Multicore Architectures (Magma) project¹⁰ is developing adaptive, dense linear algebra libraries that exploit GPUs and CPUs simultaneously. Magma is similar to Lapack, but includes support for heterogeneous architectures.

For performance and correctness tools, Georgia Tech's Ocelot project¹¹ offers a modular, dynamic compilation framework for heterogeneous systems, providing various backend targets for CUDA programs and analysis modules for the PTX virtual instruction set. Ocelot currently allows CUDA programs to be executed on Nvida GPUs, AMD GPUs, and x86-CPUs without recompilation.

Furthermore, Ocelot supports the construction of a range of correctness and performance tools, such as a memory checker that detects unaligned and out of bounds memory accesses.

For system software and virtualization, Georgia Tech is developing a framework for integrating GPUs into existing infrastructures for virtualization. This infrastructure will be useful for checkpointing and migration in virtualized systems, as well as for load balancing and debugging.

Technical Assessment

Because the architectures and software for heterogeneous computing are changing rapidly, we have an ongoing effort to evaluate different architectures and software stacks. For our assessment, we've developed the Scalable Heterogeneous Computing (SHOC) Benchmark Suite.¹³ SHOC is a collection of programs designed to test the ability of GPUs and other OpenCL devices for scalable scientific computing. SHOC has benchmarks at three levels of complexity, which measure device "feeds and speeds," important scientific kernels, and portions of full applications. SHOC also includes a stability test to validate and stress new heterogeneous architectures during procurement and installation.

Application Readiness

An important aspect of Keeneland is our applications outreach and readiness activities. In particular, we're working with early adopters to ensure that their applications work well on Keeneland. Moreover, we're aggressively pursuing those applications that might make efficient use of GPUs, but whose developers haven't yet made the jump to rewriting their applications. Our team is surveying,

contacting, modeling, and, in some cases, assisting the applications teams with porting their code to this new architecture.

Science and Applications

During our acceptance testing of the KID system in November 2010, we evaluated the system's performance with various applications and kernels for functionality, performance, and stability. The KID acceptance test included tests for all system components including the CPU, GPU, interconnect, and storage. However, for the sake of brevity, we include only the relevant GPU results here. Unless otherwise noted, the benchmarks and applications were built using the Intel 11.1.073 compilers, Intel's Math Kernel Library (MKL), OpenMPI 1.4.3, and CUDA 3.2RC.

High-performance Linpack (HPL) for TOP500 has become the dominant benchmark for high-performance computing due to its ubiquity, portability, and legacy. It has several characteristics that make it perform well on a GPU-based cluster, including reliance on dense linear algebra operations—exactly the type of regular, throughput-oriented problems that GPUs excel at solving.

Figure 3 shows KID's HPL scaling results compared to the theoretical hardware peak performance and the aggregate measured double-precision, general matrix-matrix multiplication (DGEMM) performance across GPUs. For these results, we used Nvidia's implementation of HPL, version 9. In this version, each MPI task dynamically splits work between the GPU and CPU (using multithreaded MKL). The best performance on the KID node architecture requires three MPI tasks per node, where each task controls one GPU

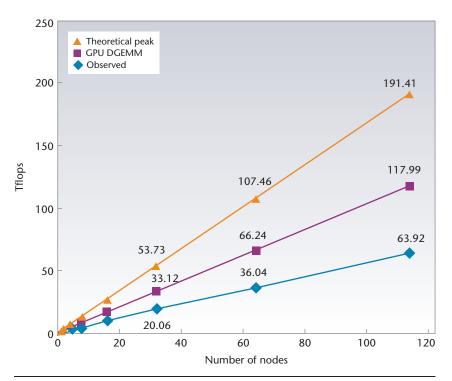


Figure 3. High-performance Linpack (HPL) performance on KID.

and uses four cores via MKL. We experienced highly variable single-node performance, presumably due to the dynamic load balancing. However, the primary bottleneck for HPL performance on KID was the relatively small size of our host memory. HPL performance is highly dependent on problem size, and nodes with a larger memory configuration are known to realize a higher percentage of peak performance. However, our configuration is sufficient for most of our real-world scientific applications.

In addition, we captured the power usage of our HPL experiments using real-time monitoring of the entire cluster. With this information, we were able to calculate the energy efficiency of KID to be 677 megaflops per watt, which placed it at the #9 ranking on the November Green500 list (www.green500.org/lists/2010/11/top/list.php). Interestingly, eight of the current top-10 systems on this list use GPUs.

Next, we ran many real-world applications on the system, including both CPU-only applications and applications that had been ported to use GPUs. Among these applications,

we found that the Groningen Machine for Chemical Simulation (Gromacs)¹⁴ and the Nanoscale Molecular Dynamics (NAMD)¹⁵ biomolecular modeling applications immediately performed well on the system without additional performance tuning. Both applications are used for biomolecular modeling. Table 2 shows performance results for these applications on KID nodes. For Gromacs, we used the dhfr-impl-2nm benchmark on a single node. We found that a single M2070 GPU outperforms a single CPU thread by 52 times, and outperforms a fully populated socket by 10.9 times. For NAMD, we used the *apoal* problem running in parallel on four nodes and realized a 6.6 times speedup when utilizing the GPUs.

Finally, during this early acceptance phase, we also ported and successfully ran the main kernel (Fast Multipole Method) for a blood flow simulation application on KID; the results were presented at the 2010 International Conference on High Performance Computing, Networking, Storage, and Analysis, where this paper was awarded the SC10 Gordon Bell prize. 16

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Table 2. Performance of a single simulation step for Gromacs and NAMD benchmarks on KID nodes.		
Benchmark	Configuration	Performance (ms)
Gromacs	1 CPU thread	39.3
	6 CPU threads	8.23
	1 GPU	0.75
NAMD	4 nodes, 12 CPU tasks	104
	4 nodes, 12 GPU tasks	15.8

urrently, many applications perform well on Keeneland, due to previous work by researchers to modify their codes to exploit GPUs. 17-19 The user community has responded positively to the availability of KID. In the first six months of operation, more than 70 projects and 200 users have requested and received access to KID. In the coming months, we'll be preparing for delivery of the final system, deploying GPU-enabled software, and engaging more applications teams so that they can use this innovative architecture to accelerate their scientific discovery, as NAMD, Gromacs, and others have.

Acknowledgments

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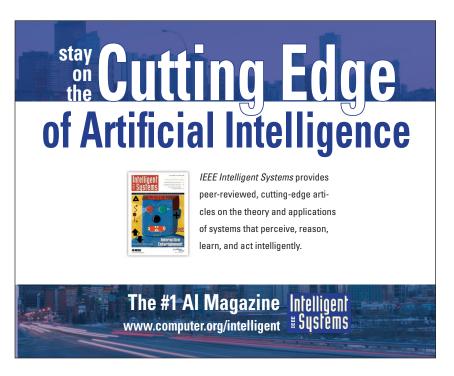
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